

Computing the Number of Cubic Runs in Standard Sturmian Words

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Run (maximal repetition)

ababaab ababaab ababaab ababaab aba ab

Notation:

- $\rho(w)$ – the number of runs in the word w
- $\rho^{(3)}(w)$ – the number of cubic runs in the word w
- $\rho(n)$ – maximal number of runs in words of length n
- $\rho^{(3)}(n)$ – maximal number of cubic runs in words of length n

a b a b a a b a b a b a b a a b a b a b a a b a b a a b
a b a b a a b a b a b a a b a b a b a a b a b a b a a b
a b a b a a b a b a b a a b a b a b a a b a b a b a a b
a b a b a a b a b a b a a b a b a b a a b a b a b a a b



Sturmian words are infinite words that has exactly $n + 1$ factors of length n for each $n \geq 0$.

Example:

a b a b a a b a b a b a a b a b a b a b a a b a b a a b ...

- ① *a b a b a a b a b a b a a b a b a b a a b a b a a b ...*
- ② *a b a b a a b a b a b a a b a b a b a a b a b a a b ...*
- ③ *a b a b a a b a b a b a a b a b a b a a b a b a a b ...*
- ④ *a b a b a a b a b a b a a b a b a b a a b a b a a b ...*



Standard word

- Alphabet: $\Sigma = \{a, b\}$
- Directive sequence: $\gamma = (\gamma_0, \gamma_1, \dots, \gamma_n)$,
where $\gamma_0 \geq 0$ and $\gamma_1, \dots, \gamma_n > 0$
- Recurrence:
 - $x_{-1} = b$
 - $x_0 = a$
 - $x_k = (x_{k-1})^{\gamma_{k-1}} \cdot x_{k-2}$
- $\text{Sw}(\gamma_0, \gamma_1, \dots, \gamma_n) = x_{n+1}$

Standard word example

- $\gamma = (1, 2, 1, 3, 1)$



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Standard word example

- $\gamma = (1, 2, 1, 3, 1)$
- $x_{-1} = b$
- $x_0 = a$
- $x_1 = (x_0)^1 \cdot x_{-1} = a \cdot b$



Standard word example

- $\gamma = (1, \textcolor{red}{2}, 1, 3, 1)$
- $x_{-1} = b$
- $x_0 = a$
- $x_1 = (x_0)^1 \cdot x_{-1} = a \cdot b$
- $x_2 = (x_1)^{\textcolor{red}{2}} \cdot x_0 = ab \cdot ab \cdot a$



Standard word example

- $\gamma = (1, 2, \textcolor{red}{1}, 3, 1)$
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- $x_3 = (x_2)^{\textcolor{red}{1}} \cdot x_1 = ababa \cdot ab$

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- $x_3 = (x_2)^1 \cdot x_1 = ababa \cdot ab$
- $x_4 = (x_3)^3 \cdot x_2 = ababaab \cdot ababaab \cdot ababaab \cdot ababa$



Standard word example

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- $x_3 = (x_2)^1 \cdot x_1 = ababa \cdot ab$
- $x_4 = (x_3)^3 \cdot x_2 = ababaab \cdot ababaab \cdot ababaab \cdot ababa$
- $x_5 = (x_4)^{\textcolor{red}{1}} \cdot x_3 = ababaabababaabababaabababa \cdot ababaab$



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- $x_5 = (x_4)^1 \cdot x_3 = ababaabababaabababaabababaabababaab$
- $Sw(1, 2, 1, 3, 1) = ababaabababaabababaabababaabababaabababaab$



Morphisms

$$\gamma = (\gamma_0, \gamma_1, \dots, \gamma_n) \quad h_i : \begin{cases} a & \longrightarrow \quad a^{\gamma_i} b \\ b & \longrightarrow \quad a \end{cases}, \quad \text{for } 0 \leq i \leq n$$



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Lemma

- ① $\text{Sw}(\gamma_n) = h_n(a)$
- ② $\text{Sw}(\gamma_i, \gamma_{i+1}, \dots, \gamma_n) = h_i\left(\text{Sw}(\gamma_{i+1}, \gamma_{i+2}, \dots, \gamma_n)\right)$



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Observation

- ① $\text{Sw}(\gamma_0, \gamma_1, \dots, \gamma_n) = h_0 \circ h_1 \circ \dots \circ h_n(a)$
- ② $\text{Sw}(\gamma_{i+1}, \gamma_{i+2}, \dots, \gamma_n) = h_i^{-1}(\text{Sw}(\gamma_i, \gamma_{i+1}, \dots, \gamma_n))$



Example

$$\gamma = (1, 2, 1, 3, 1) \quad h_i : \begin{cases} a & \longrightarrow \quad a^{\gamma_i} b \\ b & \longrightarrow \quad a \end{cases} \quad \text{for } 0 \leq i \leq n$$



Example

$$\gamma = (1, 2, 1, 3, \textcolor{red}{1}) \quad h_4 : \begin{cases} a & \longrightarrow ab \\ b & \longrightarrow a \end{cases}$$

$$S_w(1) = h_4(a) = ab$$

Example

$$\gamma = (1, 2, 1, \textcolor{red}{3}, 1) \qquad h_3 : \begin{cases} a & \longrightarrow aaab \\ b & \longrightarrow a \end{cases}$$

$$Sw(1) = h_4(a) = ab$$

$$Sw(3, 1) = h_3(Sw(1)) = aaaba$$



Example

$$\gamma = (1, 2, \textcolor{red}{1}, 3, 1) \qquad h_2 : \begin{cases} a & \longrightarrow ab \\ b & \longrightarrow a \end{cases}$$

$$Sw(1) = h_4(a) = ab$$

$$Sw(3, 1) = h_3(Sw(1)) = aaaba$$

$$Sw(1, 3, 1) = h_2(Sw(3, 1)) = abababaab$$



Example

$$\gamma = (1, \textcolor{red}{2}, 1, 3, 1) \qquad h_1 : \begin{cases} a & \longrightarrow aab \\ b & \longrightarrow a \end{cases}$$

$$Sw(1) = h_4(a) = ab$$

$$Sw(3, 1) = h_3(Sw(1)) = aaaba$$

$$Sw(1, 3, 1) = h_2(Sw(3, 1)) = abababaab$$

$$Sw(2, 1, 3, 1) = h_1(Sw(1, 3, 1)) = aabaaabaaabaaaabaaba$$



Example

$$\gamma = (1, 2, 1, 3, 1)$$

$$h_0 : \begin{cases} a & \longrightarrow ab \\ b & \longrightarrow a \end{cases}$$

$$Sw(1) = h_4(a) = ab$$

$$Sw(3, 1) = h_3(Sw(1)) = aaaba$$

$$Sw(1, 3, 1) = h_2(Sw(3, 1)) = abababaab$$

$$Sw(2, 1, 3, 1) = h_1(Sw(1, 3, 1)) = aabaaaabaaabaaaabaaba$$

$$Sw(1, 2, 1, 3, 1) = h_0(Sw(2, 1, 3, 1)) = ababaabababaababababaabababaabababaab$$



Lemma

The period of each cubic run in the standard word $S_w(\gamma_0, \dots, \gamma_n)$ is of the form x_i for $0 \leq i \leq n$.



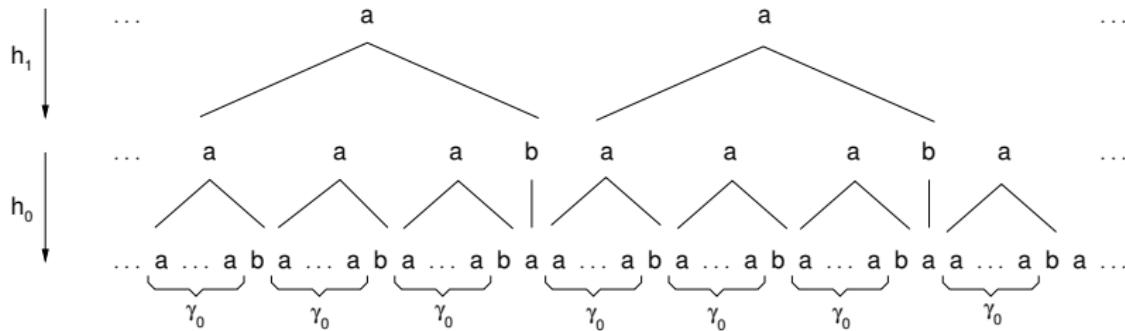
Lemma

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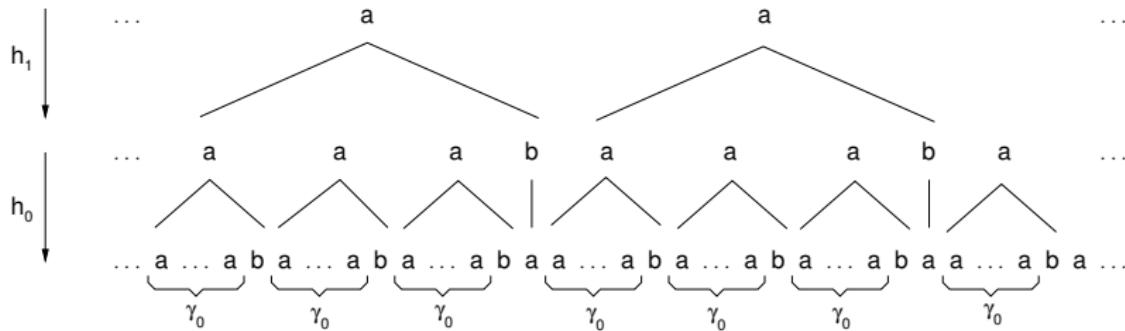
Types of cubic runs

- **short** – periods of the form x_0 and x_1 (a and $a^k b$)
- **medium** – periods of the form x_2
- **long** – periods of the form x_i for $3 \leq i \leq n$



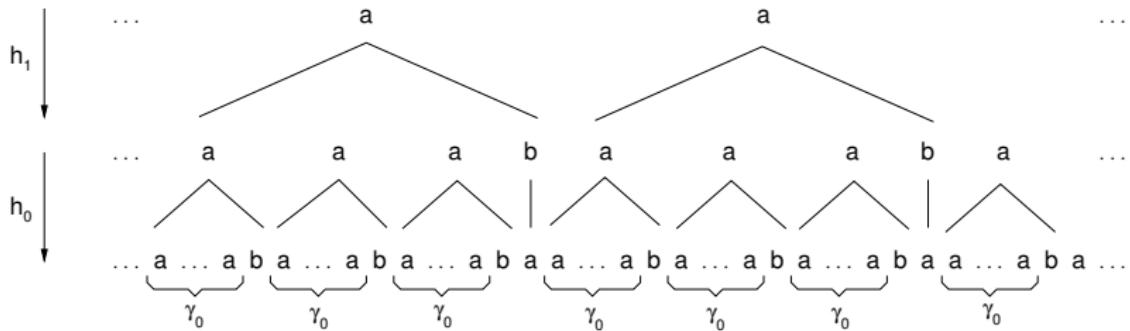


Lemma (short cubic runs with period a)



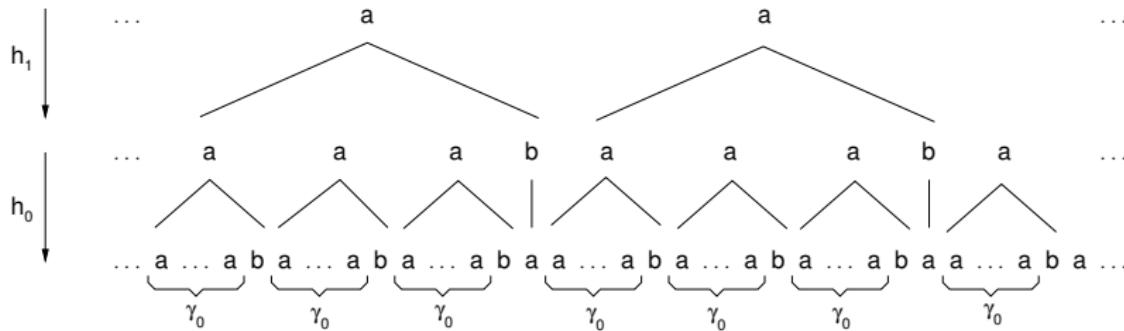
Lemma (short cubic runs with period a)

$$\rho_{S_1}^{(3)} \left(\text{Sw}(\gamma_0, \dots, \gamma_n) \right) = \begin{cases} & \\ & \left| \text{Sw}(\gamma_1, \dots, \gamma_n) \right|_a & \text{for } \gamma_0 > 2 \end{cases}$$



Lemma (short cubic runs with period a)

$$\rho_{S_1}^{(3)} \left(\text{Sw}(\gamma_0, \dots, \gamma_n) \right) = \begin{cases} \left| \text{Sw}(\gamma_2, \dots, \gamma_n) \right|_a - \text{odd}(n) & \text{for } \gamma_0 = 2 \\ \left| \text{Sw}(\gamma_1, \dots, \gamma_n) \right|_a & \text{for } \gamma_0 > 2 \end{cases}$$



Lemma (short cubic runs with period a)

$$\rho_{S_1}^{(3)} \left(\text{Sw}(\gamma_0, \dots, \gamma_n) \right) = \begin{cases} 0 & \text{for } \gamma_0 = 1 \\ \left| \text{Sw}(\gamma_2, \dots, \gamma_n) \right|_a - \text{odd}(n) & \text{for } \gamma_0 = 2 \\ \left| \text{Sw}(\gamma_1, \dots, \gamma_n) \right|_a & \text{for } \gamma_0 > 2 \end{cases}$$

Lemma (short cubic runs with period $a^k b$)

$$\rho_{S_2}^{(3)}\left(\text{Sw}(\gamma_0, \dots, \gamma_n)\right) = \begin{cases} 0 & \text{for } \gamma_1 = 1 \\ \left| \text{Sw}(\gamma_3, \dots, \gamma_n) \right|_a - \text{even}(n) & \text{for } \gamma_1 = 2 \\ \left| \text{Sw}(\gamma_2, \dots, \gamma_n) \right|_a & \text{for } \gamma_1 > 2 \end{cases}$$



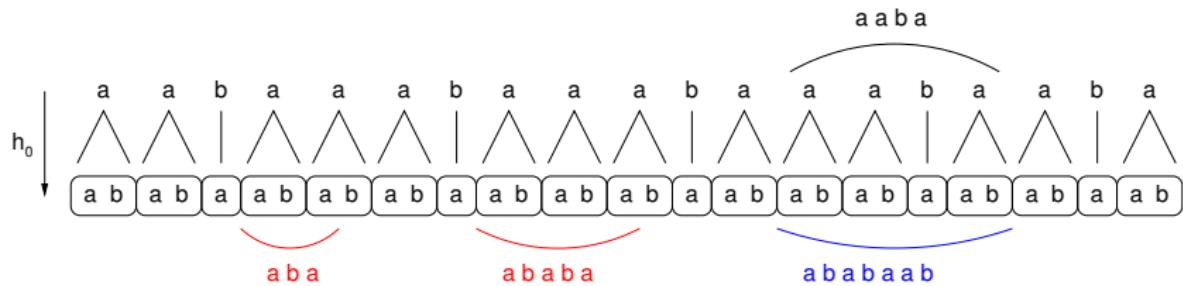
Lemma (medium cubic runs for $n \geq 3$)

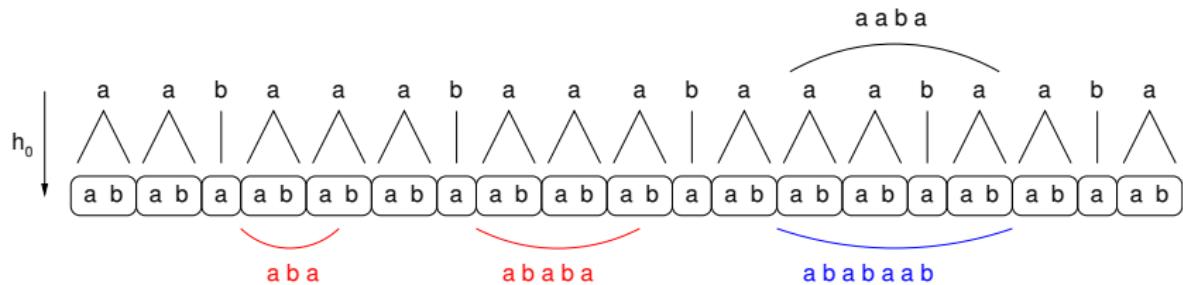
$$\rho_M^{(3)}\left(Sw(\gamma_0, \dots, \gamma_n)\right) = \begin{cases} \left|Sw(\gamma_4, \dots, \gamma_n)\right|_a - 1 & \text{for } \gamma_2 = 1 \\ \left|Sw(\gamma_3, \dots, \gamma_n)\right|_a & \text{for } \gamma_2 \geq 2 \end{cases}$$

Lemma (medium cubic runs for $n = 2$)

$$\rho_M^{(3)}\left(Sw(\gamma_0, \gamma_1, \gamma_2)\right) = \begin{cases} 0 & \text{for } \gamma_2 \leq 2 \\ 1 & \text{for } \gamma_2 > 2 \end{cases}$$







Definition

A factor x of the word w is synchronized with the morphism h iff each occurrence of x in w starts at the beginning of some h -block and ends at the end of some h -block.



Lemma

The periods of large cubic runs in the standard word $S_w(\gamma_0, \dots, \gamma_n)$ are synchronized with the morphism h_0 .



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Lemma

$$\rho_L^{(3)}(w) = \rho_M^{(3)}\left(h^{-1}(w)\right) + \rho_L^{(3)}\left(h^{-1}(w)\right)$$



Theorem

Let $w = \text{Sw}(\gamma_0, \dots, \gamma_n)$ and $w_i = \text{Sw}(\gamma_i, \dots, \gamma_n)$ be standard words. Then

$$\rho^{(3)}(w) = \rho_{S_1}^{(3)}(w) + \rho_{S_1}^{(3)}(w) + \sum_{i=0}^{n-2} \rho_M^{(3)}(w_i).$$



Theorem

We can count the number $\rho^{(3)}\left(Sw(\gamma_0, \dots, \gamma_n)\right)$ in time linear with respect to $|(\gamma_0, \dots, \gamma_n)|$, hence logarithmic time with respect to $|Sw(\gamma_0, \dots, \gamma_n)|$.



Theorem

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Observation

$$\left| \text{Sw}(\gamma_i, \dots, \gamma_n) \right|_a = \gamma_i \cdot \left| \text{Sw}(\gamma_{i+1}, \dots, \gamma_n) \right|_a + \left| \text{Sw}(\gamma_{i+2}, \dots, \gamma_n) \right|_a$$



Crochemore, Iliopoulos, Kubica, Radoszewski, Rytter, Waleń 2010

$$0.41 n \leq \rho^{(3)}(n) \leq 0.5 n$$



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$$\frac{\rho^{(3)}(F_n)}{|F_n|} \longrightarrow 0.763932 \dots$$



Theorem

For any standard word w we have:

$$\rho^{(3)}(w) \leq 0.36924841 |w|,$$

and there is strictly growing sequence $\{w_k\}$ of standard words such that

$$\lim_{k \rightarrow \infty} \frac{\rho^{(3)}(w_k)}{|w_k|} = \frac{5\Phi + 3}{13\Phi + 9} \approx 0.36924841,$$

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Idea of the proof

$$w_k = (3, 2, \underbrace{1, 1, \dots, 1}_k)$$



Thank You
For Your Attention

